

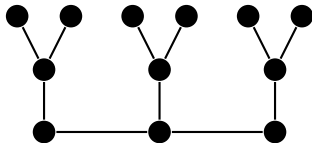
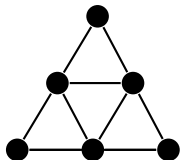
Fun with $L(2,1)$ -labeling

Brian Heinold

Department of Mathematics and Computer Science
Mount St. Mary's University

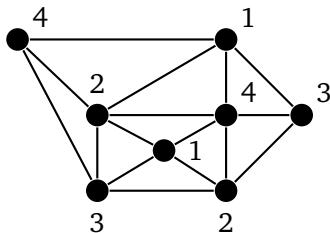
April 29, 2023

Network of points (vertices) and lines (edges)



Graph coloring

Assign colors (usually numbers) to vertices so that neighbors get different colors. Try to use as few colors as possible.



Assign numerical labels to vertices according to the following rules:

- 1 Neighbors' labels must differ by 2.
- 2 Vertices at distance 2 must get different labels.

Start with 1 and try to keep the labels as small as possible.

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Breaks rule #1

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Breaks rule #2

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Okay, but not optimal.

L(2,1)-labeling

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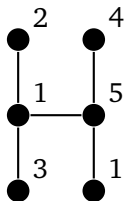
Optimal!

Another example

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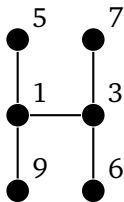
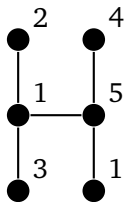


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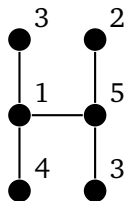
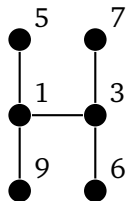
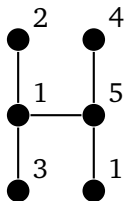


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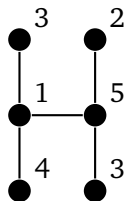
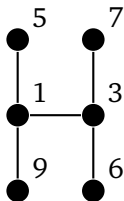
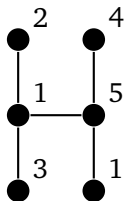


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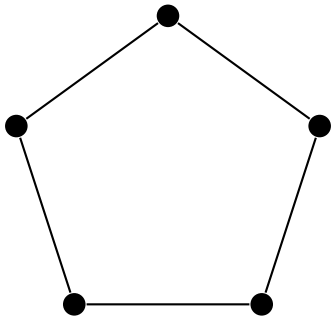
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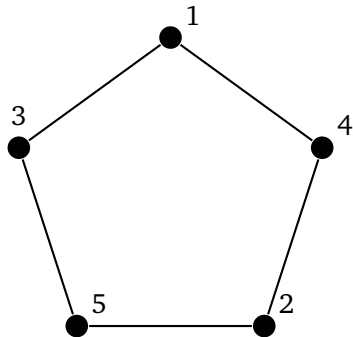


Left graph breaks both rules repeatedly. Middle graph is okay but uses too many colors. Right graph is optimal.

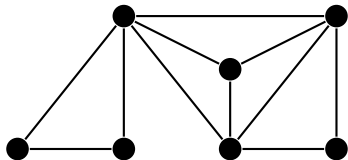
Let's try another



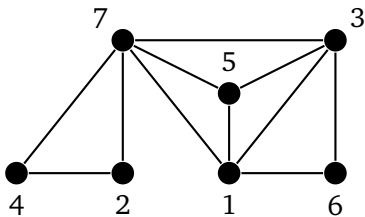
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Yet another example

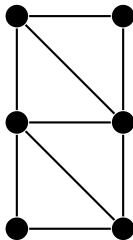


Yet another example



A challenge

Goal: max label 6



About the problem

- Introduced in 1992 by Griggs and Yeh, based on an idea of F. Roberts.

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- People are interested in minimizing the largest label used.
- We'll denote it by $\lambda_{2,1}(G)$.
- In the literature, people often talk about the *span*, the difference between the largest and smallest labels.

- $\lambda_{2,1}$ can be computed quickly for various classes of graphs like paths, cycles, complete graphs, etc. These are good exercises.

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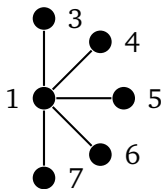
- $\lambda_{2,1}$ can be computed quickly for various classes of graphs like paths, cycles, complete graphs, etc. These are good exercises.
- There are many classes of graphs for which $\lambda_{2,1}$ is not known.
- The problem is NP-complete (i.e., interesting).

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Lower bound

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- See below for an example.

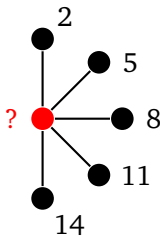


Upper bound

- With a little more work, we can show the max label $\lambda_{2,1}$ is always less than $\Delta^2 + 2\Delta + 1$.

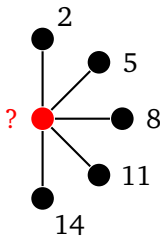
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- Why? A worst-case scenario is shown below. We're trying label the middle left vertex. Assume its neighbors have already been labeled. In this worst-case scenario, we have to avoid labels 1 through 15. In general, this could be 3Δ labels.



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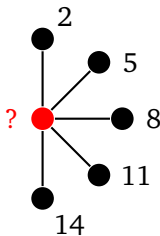
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- In total, we have to avoid at most $3\Delta + \Delta(\Delta - 1) = \Delta^2 + 2\Delta$ labels.

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- Sufficiently large here means 10^{69} .
- They also showed that $\lambda_{2,1} \leq \Delta^2 + c$ for some fixed (but unfortunately very large) c .

Relatives and generalizations

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- Can also look at $L(p, q, r)$, $L(p, q, r, s)$, etc.
- There are also edge versions, list versions, and more.

- MR4556054** Pending Zhang, Xiaoling Optimal $L(2,1,1)$ -labelings of caterpillars. *J. Math. Res. Appl.* 43 (2023), no. 2, 150–160. [05C78](#)
[Review PDF](#) | [Clipboard](#) | [Journal](#) | [Article](#)
- MR4568901** Prelim Amanathulla, Sk.; Bera, Biswajit; Pal, Madhumangal; $L(2,1,1)$ -labeling of interval graphs. *Int. J. Math. Ind.* 14 (2022), Paper No. 2250003, 8 pp. [05C78](#)
[Review PDF](#) | [Clipboard](#) | [Journal](#) | [Article](#)
- MR4496569** Reviewed Zhu, Haiyang; Zhu, Junlei; Liu, Ying; Wang, Shuling; Huang, Danjun; Miao, Lianying Optimal frequency assignment and planar list $L(2,1)$ -labeling. *J. Comb. Optim.* 44 (2022), no. 4, 2748–2761. [05C10](#) ([05C78](#))
[Review PDF](#) | [Clipboard](#) | [Journal](#) | [Article](#)
- MR4345146** Reviewed Colucci, Lucas; Györi, Ervin On $L(2,1)$ -labelings of oriented graphs. *Discuss. Math. Graph Theory* 42 (2022), no. 1, 39–46. [05C78](#) ([05C20](#))
[Review PDF](#) | [Clipboard](#) | [Journal](#) | [Article](#)
- MR4387595** Reviewed Ma, Dengju; Yao, Shunyu; Dong, Xiaoyuan The $L(3,2,1)$ -labeling number of the Cartesian product of a complete graph and a cycle. *Ars Combin.* 154 (2021), 87–100. [05C78](#)
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- MR4387590** Reviewed Lv, Damei; Lin, Nianfeng $L(2,1)$ -circular-labelings of the edge-path-replacements. *Ars Combin.* 154 (2021), 23–30. [05C78](#)
[Review PDF](#) | [Clipboard](#) | [Journal](#) | [Article](#)
- MR4299390** Reviewed Bandopadhyay, Susobhan; Ghosh, Sasthi C.; Koley, Subhasis Improved bounds on the span of $L(1,2)$ -edge labeling of some infinite regular grids. *Graphs and combinatorial optimization: from theory to applications—CTW2020 proceedings*, 53–65, *AIRO Springer Ser.*, 5, Springer, Cham, [2021], ©2021. [05C78](#) ([05C12](#))
[Review PDF](#) | [Clipboard](#) | [Series](#) | [Chapter](#)

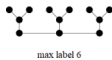
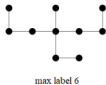
My assignment

Name _____

Math 211 Assignment 33 (Due Fri 12/2/22)

$L(2, 1)$ -labeling is a relative of graph coloring. Vertices are assigned integer labels according to the following rules:

- The labels for neighbors must differ by at least 2.
- If a vertex is two steps away from another (i.e., to get to it, you have to pass through one exactly other vertex), then their labels need to be different.
- If vertices are more than two steps away from each other, then it is okay for them to have the same label.
- The "max label" listed under each graph is the largest number you are allowed to use. That label should show up somewhere in the graph. The smallest label used should always be 1. It's possible that not all the numbers between 1 and the max will be needed.



- I've used it for several years. Students in a non-majors course have told me they found it
 - Fun
 - Makes their brain hurt
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 - Developing a habit of checking work for mistakes
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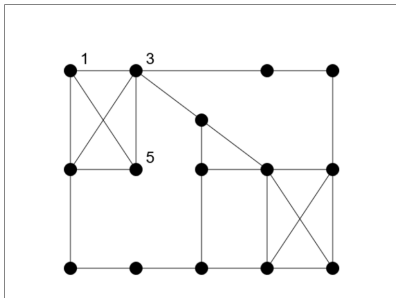
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 - Developing perseverance, since first attempts often don't work
 - Developing a habit of checking work for mistakes
 - Developing problem-solving strategies
- It's not something they have ever seen before.
- It's also a good problem for undergraduate research since it doesn't require a lot of background to get started on, and there are myriad things to work on.

- I find $L(2, 1)$ -labelling works well as a puzzle. It's easy to make mistakes, though. I built a simple app that checks your work.

<https://www.brianheinold.net/L21.html>

- You can view source on the page to see the code. Adding new graphs is quick.

Board number:

Goal: 7

Rules:

1. Neighbors' labels must differ by 2.
2. Vertices at distance 2 from each other must get different labels.
3. The smallest label is 1 and the largest label shouldn't be more than the goal.

Thank you!

A few selected references:

T. Calamoneri. The $L(h, k)$ -Labeling Problem: A Survey, (2014)

<http://www.dsi.uniroma1.it/calamo/PDF-FILES/survey.pdf>

J.R. Griggs and R.K. Yeh, Labeling graphs with a condition at distance two, SIAM J. Discr. Math. 11 (1992), 585–595.

Havet, F., Reed, B. and Sereni, J.-S. (2008) $L(2, 1)$ -Labeling of graphs. Proceedings of the ACM-SIAM Symposium on Discrete Algorithm (SODA '08), San Francisco, California, 20-22 January, 621–630.